Bitcoin & Cryptocurrency

Material taken from <u>here</u>, <u>here</u>

- Alice wants to digitally send some money to Bob
- Need identities
 - Each user has a public-private key pair
 - Each user referred to by their public key



- Alice simply signs a message saying that she is giving Bob a coin
- Good properties:
 - Publicly verifiable
 - Alice cannot renege
 - No one can forge this transaction



"I, Alice, am giving Bob one coin", σ

- Problem: what if Alice keeps on sending the same message over and over again, even if she doesn't have enough coins?
 - Double spending
- Coins need to be unique
- Banks can handle this as a centralized authority using a ledger





- "I, Alice, am giving Bob one coin", σ
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- Ledger keeps track of initial budgets, and all transactions
- Ledger owner checks that
 - Signature verifies for sender with *PK*
 - Transaction is indeed sending from a sender with the corresponding *PK*
 - Checks for any double spending

$$PK_A$$
: 5 coins $TX_1 = (PK_A \rightarrow PK_B,$ PK_B : 0 coins $\sigma_{SK_A}(TX_1)$



"I, Alice, am giving Bob one coin", σ

1): 2 coins	$TX_2 = (PK_A \rightarrow PK_B, 1): 5 \text{ co}$ $\sigma_{SK_A}(TX_2)$



Making everyone collectively the bank

- Everyone keeps a copy of the ledger using a blockchain
 - Each transaction contains a hash of the previous transaction
 - Given h(block i) from a trusted source, and preceding blocks from an untrusted source, can verify that the previous blocks are not compromised
- Updates need to be propagated to everyone
- Everyone checks the transaction, and if it's correct, creates a new block including this transaction and add to the local blockchain

$$PK_A$$
: 5 coins $TX_1 = (PK_A \rightarrow PK_B,$ PK_B : 0 coins $\sigma_{SK_A}(TX_1)$ $h(block1)$

1): 2 coins	$TX_2 = (PK_A \rightarrow PK_B, 1)$: 5 co
	$\sigma_{SK_A}(TX_2)$
	h(block2)



Making everyone collectively the bank

- Problem: Alice purchases an item from Bob for 100 coins. Then she create new blocks of transactions.
- Solution: proof-of-work

 - Reward them for trying to help validate transactions

creates a fork from the block just before the transaction, and she can

Make it computational costly for network users to validate transactions

Proof-of-work

- Only miners can add blocks to the blockchain
- All miners try to solve a computationally hard (but solvable) problem: find a nonce such that h(block) starts with a number of zeros
 - Bitcoin protocol actually requires the hash to be less than or equal to a <u>"target"</u>: every 2016 blocks, every Bitcoin client compares the actual time it took to generate these blocks and modifies the target
- Once a miner solves a proof-of-work, it creates a block that includes a set of received transactions
- Miner receives a free coin if mining is successful

Consensus

- Longest correct chain wins; everyone checks all blocks and all transactions
- Can Alice create a fork?
 - Bob waits for a number of blocks to be appended after the said transaction. If Alice does not have a majority of the computing power, then the probability that she can catch up drops exponentially.
- What happens if Alice and Bob are mining blocks, and both solve a proofof-work and append two different blocks, creating a fork?
 - The next miner that appends onto one of these chains invalidates the other chain.

Consensus

- If a miner has included Alice's transaction in the latest block, is it guaranteed that her transaction is forever in the blockchain?
 - No, there could be a legitimate miner appending a different block, so it's safe to wait for a few blocks.
- What happens if a miner refuses to include Alice's transaction?
 - It's unfortunate, but there could be another miner that will do so. Each transaction includes a fee that goes to the miner so a miner is incentivized to include as many transactions as possible.

Bitcoin is anonymous?

- Silk Road: online black market
- SoK on this topic
 - Common input ownership
 - Change addresses
 - Transaction graph analysis
 - Side-channel information

Ethereum

Next time: guest lecture!

Title: Proof-of-stake blockchains and compact certificates

Abstract: In this lecture I will sketch the big ideas behind proof-of-stake blockchains, using Algorand as the running example. Then I will define a cryptographic primitive called a compact certificate and discuss how compact certificates can be used to take a snapshot of a proof-of-stake blockchain. At the end, I will show how these snapshots can be used both to help clients quickly catch up to the current state of the chain, and to implement cross-blockchain interoperation.